APT Session 5: Interpreters



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Software Development Team 2015-01-27

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What to expect from this session

- 1 How do programming languages run programs?
- 2 Building your own interpreter.

Prerequisites

1 Have the programming language of your choice (e.g. Java, Python) installed and running on your computer.

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- *1* Compiling down into machine code at compile-time (e.g. C).
- 2 Compiling to machine code at run-time (e.g. Java).
- *3* Having another program *interpret* your program at run-time (e.g. Python).

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- Although interpreters aren't particularly fast (on their own), they're fast enough that they're used heavily in the real-world.

Our language

We're going to build an interpreter for a simple stack-based language. Here's an example program:

INT 2 Push 2 onto the	stack
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- INT 3 Push 3 onto the stack
- ADD Pop the last 2 elements from the stack, add them, and push the result onto the stack
- PRINT Peek the last element from the stack and print it

Terminology:

Stack A FILO (First In Last Out) list.

- *Push* Add an element to the top of the stack.
 - *Pop* Remove the top-most element from the stack for inspection.
- *Peek* Inspect the top-most element of the stack & don't remove it.

Parsing

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- *1* Put the above program into a file pl.myl.
- 2 Write a program which reads the file in and splits each instruction into a list of strings. You may assume that every instruction has a name and 0 or 1 arguments. The list in memory should look roughly like:

```
[["INT", "2"], ["INT", "3"], ["ADD", ""], ["PRINT", ""]]
```

A basic interpreter

We now have a list in memory along the lines of:

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Exercises:

- 1 Write the main loop of the interpreter, specifying only ADD and PRINT instructions. I suggest that all elements on the stack are stored as integers.
- 2 Run pl.myl

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Control flow

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- ¹ Implement the SUB instruction: it pops (in order) elements e_1 and e_2 , performs $e_2 e_1$ and puts it back on the stack. [NB: We didn't need to be this careful for ADD because addition is *commutative*.]
- 2 Implement the JGE x instruction. It peeks at the top-most element: if it is ≥ 0 it jumps to the instruction at position x; otherwise it adds 1 to the PC.
- 3 Store the program above as p2.myl and run it.

Procedures

Jumps can build for/while loops (etc.) but not function/procedures.

1:	INT 100	6:	INT 1
2:	CALL 4	7:	SUB
3:	JGE 1	8:	PRINT
4:	EXIT	9:	SWAP
5:	SWAP	10:	RET

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- *1* Implement the SWAP instruction which swaps the two top-most elements on the stack around.
- 2 Implement the CALL *x* instruction: it pushes the pc + 1 onto the stack and jumps to position *x*.
- *3* Implement the RET instruction: it pops the top-most value from the stack and jumps to that value.
- 4 Implement the EXIT instruction: it exits the program.
- 5 Store the program above as p3.myl and run it.

Labels

Jumping to numeric offsets is fragile. Labels make programs more robust:

- INT 100
- L1: PRINT
 - INT 1
 - SUB
 - JGE L1

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- 1 Allow users to define labels before an instruction and to jump to it later. [NB: Labels can be defined *after* a jump which references them.]
- 2 Store the program above as p4.myl and run it.

Fibonacci

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- 1 Write the Fibonacci program in your language. You will probably need to add DUP (peek at the top-most element on the stack and push a copy of it), JEQ x (peek at the top-most element of the stack and if it is 0 jump to pc x), and POP (discard the top-most element of the stack).
- 2 Store the program above as fib.myl and run it.

Try these (no particular order):

- Convert your interpreter to use integer constants instead of strings to represent instructions (tends to give a small speed-up).
- Rewrite your interpreter in RPython and have a working JIT!