

Automatically Retrofitting JIT Compilers

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yk

- *VM:*

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- *JIT compiler*:

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- *VM*: System containing one or more language implementations
- *Interpreter*: Simple language implementation*
- *JIT compiler*: Language implementation that observes a running program, optimises, and compiles portions to machine code

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I: Why?

My program is slow

Drop in a faster language implementation

Cinder; IronPython; Jython;

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Nuitka; Psyco; Pyjion;

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Shed Skin; Stackless; Starkiller;

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PyPy; Pyston; S6;
Shed Skin; Stackless; Starkiller;
TrufflePython; Unladen Swallow;

Cinder; IronPython; Jython;
Nuitka; Psyco; Pyjion;
PyPy; Pyston; S6;
Shed Skin; Stackless; Starkiller;
TrufflePython; Unladen Swallow;
WPython; Zippy

JIT compiling VMs are:

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hard

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expensive

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maybe incompatible

JIT compiling VMs are:

hard

expensive

maybe incompatible

difficult to evolve

Automation:

Automation: RPython and Truffle

Automation: RPython and Truffle

Host:

Automation: RPython and Truffle

Host: RPython / Java

Automation: RPython and Truffle

Host: RPython / Java

Guest:

Automation: RPython and Truffle

Host: RPython / Java

Guest: Lua, Ruby, Python, CPU sim, etc.

II: How?

The C interpreter is the source of truth

Generate a *meta-tracing* JIT compiler
from C interpreters

Tracing: manually record *hot loops* at run-time

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Meta-tracing: record the interpreter
executing loops at run-time

AOT



Run-time

AOT



Run-time

AOT



ykllvm



Run-time

AOT



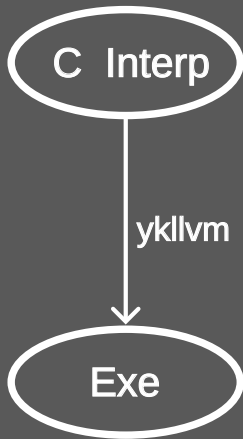
ykllvm



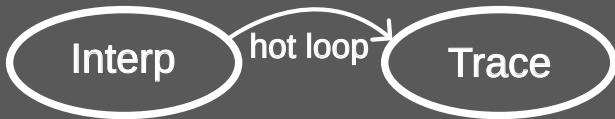
Run-time



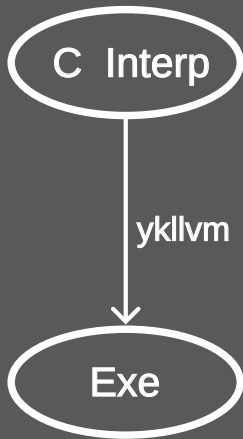
AOT



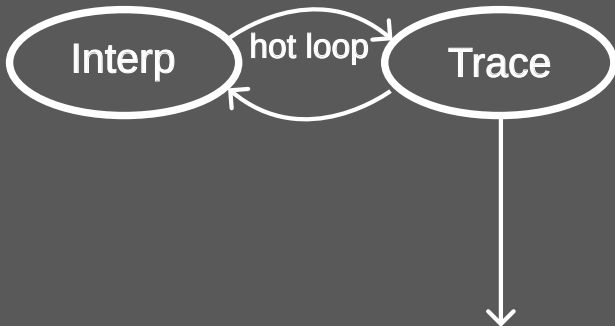
Run-time



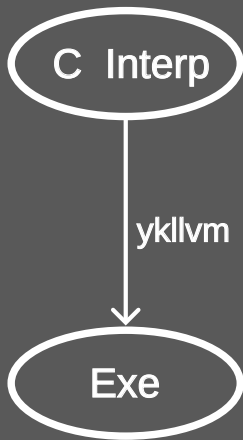
AOT



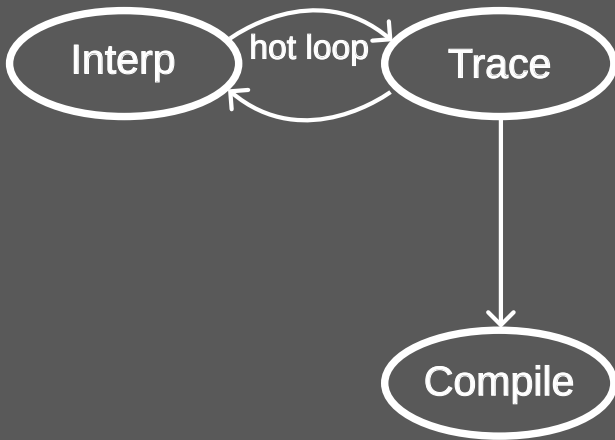
Run-time



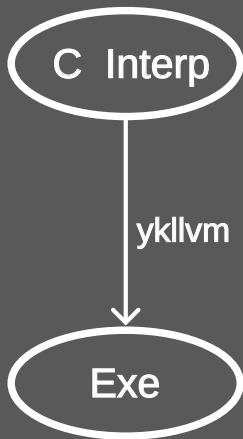
AOT



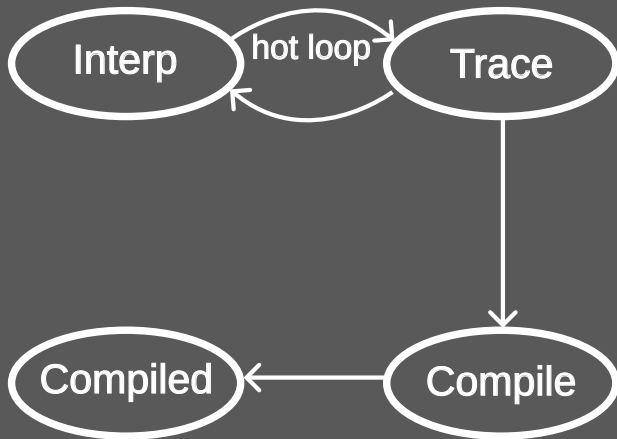
Run-time



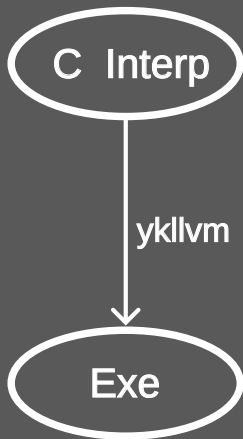
AOT



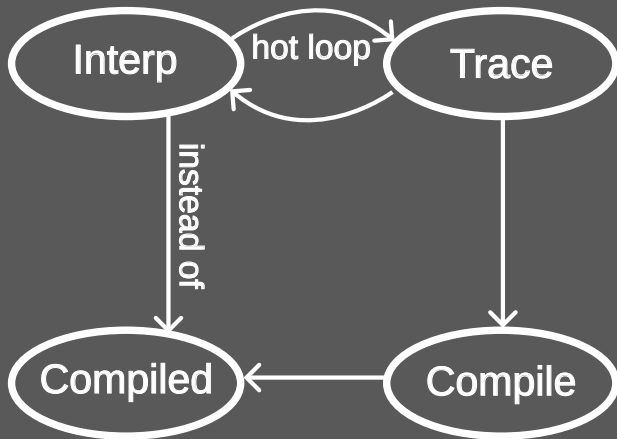
Run-time



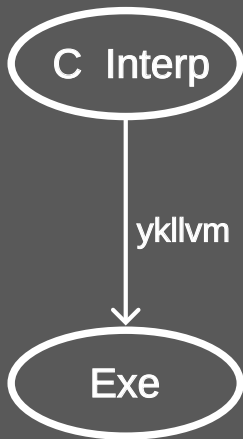
AOT



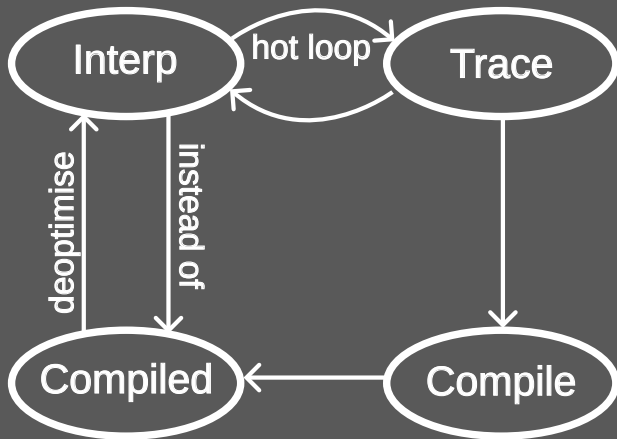
Run-time



AOT



Run-time



```
while (true) {
```

```
}
```

```
Instruction *code = ...;
```

```
while (true) {
```

```
}
```

```
Instruction *code = ...;
int pc = 0;

while (true) {
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {

        }
    }
}
```

```
Instruction *code = ...;
int pc = 0;

while (true) {
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP:

    }
}
```

```
Instruction *code = ...;
int pc = 0;
Value **stack = ...;
while (true) {
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP: push(lookup(GET_OPVAL()));

    }
}
```



```
Instruction *code = ...;
int pc = 0;
Value **stack = ...;
while (true) {
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP: push(lookup(GET_OPVAL())); pc++; break;

    }
}
```

```
Instruction *code = ...;
int pc = 0;
Value **stack = ...;
while (true) {
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP: push(lookup(GET_OPVAL())); pc++; break;
        case OP_ADD:

    }
}
```

```
Instruction *code = ...;
int pc = 0;
Value **stack = ...;
while (true) {
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP: push(lookup(GET_OPVAL())); pc++; break;
        case OP_ADD: push(pop() + pop()); pc++; break;

    }
}
```

```
Instruction *code = ...;
int pc = 0;
Value **stack = ...;
while (true) {
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP: push(lookup(GET_OPVAL())); pc++; break;
        case OP_ADD: push(pop() + pop()); pc++; break;
        case OP_JLE:

    }
}
```

```
Instruction *code = ...;
int pc = 0;
Value **stack = ...;
while (true) {
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP: push(lookup(GET_OPVAL())); pc++; break;
        case OP_ADD: push(pop() + pop()); pc++; break;
        case OP_JLE:
            if (pop() <= 0) pc += GET_OPVAL(i) else pc++;
            break;
    }
}
```


Guest

```
if x > 0:  
    y = y + 3
```

Tracing

```
OP_LOOKUP("x")  
guard true, pop() > 0  
OP_LOOKUP("y")  
OP_INT(3)  
OP_ADD  
OP_SET("y")
```

Guest

```
if x > 0:  
    y = y + 3
```

Tracing

```
OP_LOOKUP("x")  
guard true, pop() > 0  
OP_LOOKUP("y")  
OP_INT(3)  
OP_ADD  
OP_SET("y")
```

Meta-tracing

```
push(lookup("x")); pc++;  
guard true, pop() > 0; pc++;  
push(lookup("y")); pc++;  
push(3); pc++;  
push(pop() + pop()); pc++;  
set("y", pop()); pc++;
```


How does yk convert C code into a meta-tracing compiler?

How does yk convert C code into a meta-tracing compiler?

yklvm


```
Instruction *code = ...;
int pc = 0;
Value **stack = ...;
while (true) {

    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP:

            push(lookup(GET_OPVAL())); pc++; break;
        case OP_ADD:

            push(pop() + pop()); pc++; break;
    }
}
```

```
Instruction *code = ...;
int pc = 0;
Value **stack = ...;
while (true) {
    __yk_record(0);
    Instruction i = code[pc];
    switch (GET_OPCODE(i)) {
        case OP_LOOKUP:
            __yk_record(1);
            push(lookup(GET_OPVAL())); pc++; break;
        case OP_ADD:
            __yk_record(2);
            push(pop() + pop()); pc++; break;
    }
}
```


Convert LLVM's IR for the interpreter into yk IR

Executable = normal machine code + serialised IR

How does yk optimise a program?

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1 Inlining.

How does yk optimise a program?

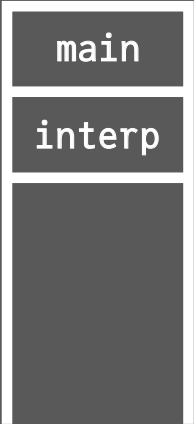
- 1 Inlining.
- 2 Standard(ish) compiler optimisations.

How does yk optimise a program?

- 1 Inlining.
- 2 Standard(ish) compiler optimisations.
- 3 Interpreter hints.

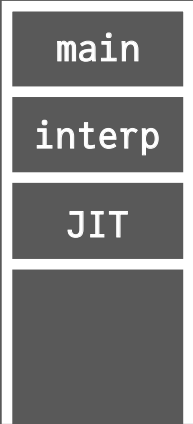


main



main

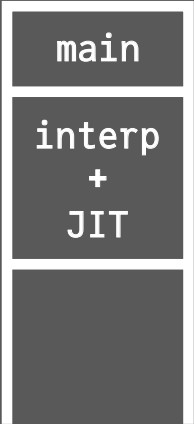
interp



main

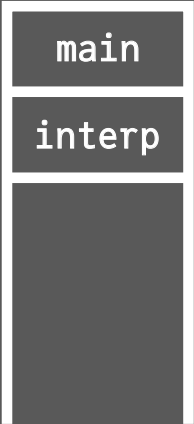
interp

JIT



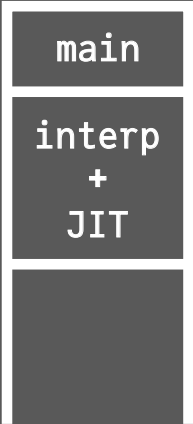
main

interp
+
JIT



main

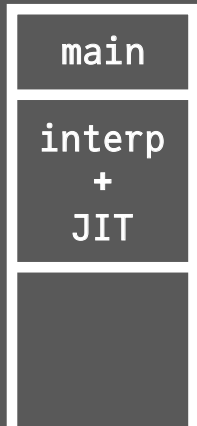
interp



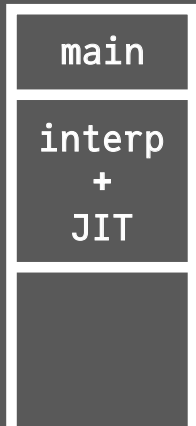
main

interp
+
JIT

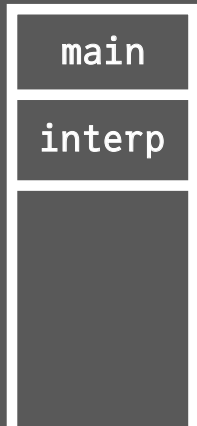
"C"



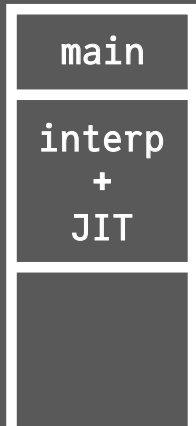
Shadow



"C"



Shadow




```
while i > 0 do  
    x = ...  
    i = i - 1  
end  
print(x)
```

```
while i > 0 do  
    x = ...  
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```

Is this a loop?

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while i > 0 do  
    x = ...  
    i = i - 1  
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Is this a loop?

AOT:

```
while i > 0 do  
    x = ...  
    i = i - 1  
end  
print(x)
```

Is this a loop?

AOT: yes

```
while i > 0 do  
    x = ...  
    i = i - 1  
end  
print(x)
```

Is this a loop?

AOT: yes

Run-time:


```
while i > 0 do  
    x = ...  
    i = i - 1  
end  
print(x)
```

Is this a loop?

AOT: yes

Run-time: maybe

```
while i > 0 do  
    x = ...  
    i = i - 1  
end  
print(x)
```

```
start: LOOKUP("i")  
      JLE(end)  
      ...  
      SET("x")  
      LOOKUP("i")  
      SUBI(-1)  
      JMP(start)  
end:  LOOKUP("x")  
      PRINT()
```

```
while i > 0 do  
    x = ...  
    i = i - 1  
end  
print(x)
```

```
start: LOOKUP("i")  
      JLE(end)  
      ...  
      SET("x")  
      LOOKUP("i")  
      SUBI(-1)  
      JMP(start)  
end:   LOOKUP("x")  
      PRINT()
```



```
while i > 0 do  
    x = ...  
    i = i - 1  
end  
print(x)
```

```
start: LOOKUP("i")  
      JLE(end)  
      ...  
      SET("x")  
      LOOKUP("i")  
      SUBI(-1)  
      JMP(start)  
end:   LOOKUP("x")  
      PRINT()
```



The diagram illustrates a loop structure. A curved arrow originates from the 'end' label and points back to the 'start' label, indicating a jump back to the beginning of the loop body. Another curved arrow originates from the 'JLE(end)' instruction and points to the '...' instruction, indicating a jump back to the start of the loop body.

III: What's next?

1 Support more LLVM IR (e.g. vectors).

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- 2 More optimisations (e.g. escape analysis).

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- 2 More optimisations (e.g. escape analysis).
- 3 More interpreters (WIP micropython).

<https://github.com/ykjit/yk>